

University of Bahrain
College of Information Technology
Department of Computer Science
Semester 2, 2012-2013

ITCS312/ITCS314 (Formal languages and Automata Theory)

Test-II

Date: May 13, 2013

Time: 3:00-4:00

STUDENT NAME	DRAGON
STUDENT ID #	
SECTION	

NOTES:

- WRITE ONLY ONE SOLUTION FOR EACH QUESTION.
- SWITCH OFF YOUR MOBILE PHONES.
- THIS EXAM CONTAINS 4 PAGES INCLUDING THE COVER PAGE.
- ANSWER ALL THE FOLLOWING QUESTIONS
- $\Sigma = \{a,b\}$ UNLESS STATED OTHERWISE

QUESTION #	MARKS		COMMENTS
1	8	6	
2	12	10 ³ / ₄	
3	10	9 ³ / ₄	
TOTAL	30	26 ¹ / ₂	

Question One (True/ False) (8 marks)

1. $a^n b^{2n}$ is a non-regular language (True). ✓
2. \wedge will be in Pref(Q in R) if there is no common words in Q and R. (False) ✓
3. The quotient of two regular languages is unique. (False) ✓
4. If number of states in FA is 4 and Σ has 3 alphabets, then the number of the possible accepted words between n and less than 2n in length is 243. (False) ✓
5. A CFG is called a RG if the language productions of the form (False) ✓
 Nonterminal \rightarrow working string
 Nonterminal \rightarrow word
6. If $R1 = R2$, then the FA representing the intersection has always accepted words. (True) ✓
7. The total tree can be sometimes finite even if we have infinite language. (False) ✓
8. The language is ambiguous if it can be represented by more than one parse tree (True) ✓

6

Question Two (12 marks)

1. Give the CFG for the language $a^n b^{2n}$ where $n \geq 0$. (2 marks)

$S \rightarrow aSbb \mid \lambda$ ✓ 2

2. Give CFG for $a^x b^y a^z$ where $z = x - y$ & $x, y, z \geq 0$ (2 marks)

$S \rightarrow \lambda \mid aXa$
 $X \rightarrow aXa \mid aRba \mid \lambda$
 $R \rightarrow \lambda$

1 how you get
 $aaabba$ for example

3. Given the following CFG, give another CFG to show that this language is regular. (3 marks)

$S \rightarrow aaXbb$

$X \rightarrow aX \mid bX \mid a \mid b$

$aa(a+b)(a+b)^*bb$

$S \rightarrow aaR_1$
 $R_1 \rightarrow aR_2 \mid bR_2$
 $R_2 \rightarrow aR_2 \mid bR_2 \mid bb$ ✓ 3

10^{3/4}

4. Give the CFG for $(a+bb)abba(ab+aa)$ (3 marks)

$$S \rightarrow XabbaY$$

$$X \rightarrow a|bb$$

$$Y \rightarrow ab+aa$$

5. Given the following CFG, use leftmost derivation to derive **aaabb** if derivable.
(2 marks)

$$S \rightarrow XY$$

$$X \rightarrow XX | a$$

$$Y \rightarrow YY | b$$

$$S \Rightarrow XY \Rightarrow XX Y \Rightarrow aXY \Rightarrow aXXY \Rightarrow aaXY \Rightarrow$$

$$\Rightarrow aaaY \Rightarrow aaaYY \Rightarrow aaabY \Rightarrow aaabb$$

Question 3 (10 marks)

1. Convert the following CFG to CNF (3 marks)

$$S \rightarrow Ab | bb | ABa$$

$$A \rightarrow \wedge | BR | a$$

$$B \rightarrow b | RMA | aa$$

$$R \rightarrow b | \wedge | S$$

$$M \rightarrow aba | \wedge$$

A, R, M are Nullables

stage 1: remove \wedge productions

$$S \rightarrow Ab|bb|ABa|b|Ba|aM \rightarrow A_1B_1A_1$$

$$A \rightarrow BR|a|B|R$$

$$B \rightarrow b|RMA|aa|MA|RM|RA|RM|A$$

$$R \rightarrow b$$

$$M \rightarrow aba$$

stage 2:

$$S \rightarrow AB_1|B_1B_1|ABA_1|b|BA_1$$

$$A \rightarrow BR|a|b|RMA|A_1A_1|MA|RM|RA|b|$$

$$A_1B_1A_1|BR|a$$

$$B \rightarrow b|RMA|A_1A_1|MA|RM|RA|b|A_1B_1A_1|$$

$$BR|a$$

$$R \rightarrow b$$

stage 3: $S \rightarrow AB_1|B_1B_1|AR_1|b|BA_1$

$$A \rightarrow BR|a|b|RR_2|A_1A_1|MA|$$

$$RM|RA|A_1R_3|BR|a$$

$$B \rightarrow b|RR_2|A_1A_1|MA|RM|RA|A_1R_3|BR|a$$

$$R \rightarrow b$$

$$M \rightarrow A_1R_3$$

$$R_1 \rightarrow BA_1$$

$$R_2 \rightarrow MA$$

$$R_3 \rightarrow B_1A_1 \Rightarrow$$

$$A_1 \rightarrow a$$

$$B_1 \rightarrow b$$

9 3/4

2. Using Pumping lemma, verify whether the language $a^n b^m$ is regular or not, where $n \geq m$ & $n, m \geq 0$. (3 marks)

assume the language is regular and has FA with r states.

$\therefore a^r b^r$ should be in the language. (if $n=m$)

$a^r = uvw$ such that $|v| \leq r$

\therefore for $k \geq 0$, $uv^k w b^r$ is in the language

but for $k=0$, $uw b^r$ is also in the language, but uw contains less a 's than a^r , \therefore the language is not regular.

3. Check whether the following grammar ambiguous or unambiguous. (2 marks)

$S \rightarrow baX \mid bA \mid aaM$

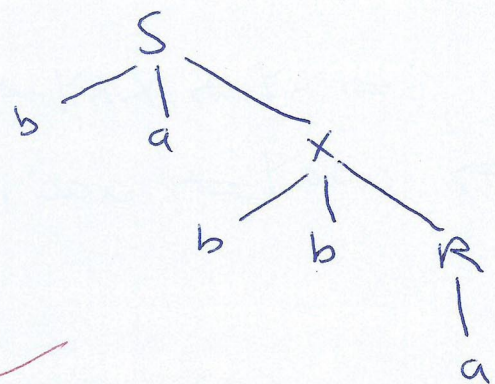
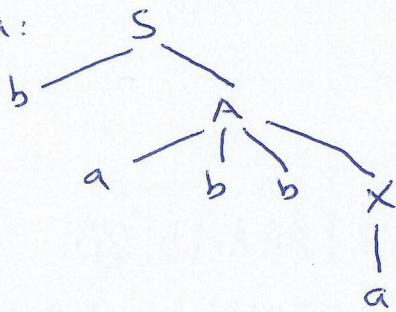
$M \rightarrow abS \mid bb$

$X \rightarrow bbR \mid a$

$A \rightarrow abbX \mid \wedge$

$R \rightarrow a \mid b$

for babba:



\therefore there are two parse trees for the same string.

\therefore the grammar is ambiguous.

4. Given the following CFG, give another CFG without unit production (2 marks)

$S \rightarrow AA$

$A \rightarrow B \mid BB$

$B \rightarrow abB \mid b \mid bb$

solution:

$S \rightarrow AA$

$A \rightarrow abB \mid b \mid bb \mid BB$

$B \rightarrow abB \mid b \mid bb$

$A \rightarrow B$

$A \Rightarrow B \Rightarrow abB$

$A \Rightarrow B \Rightarrow b$

$A \Rightarrow B \Rightarrow bb$